

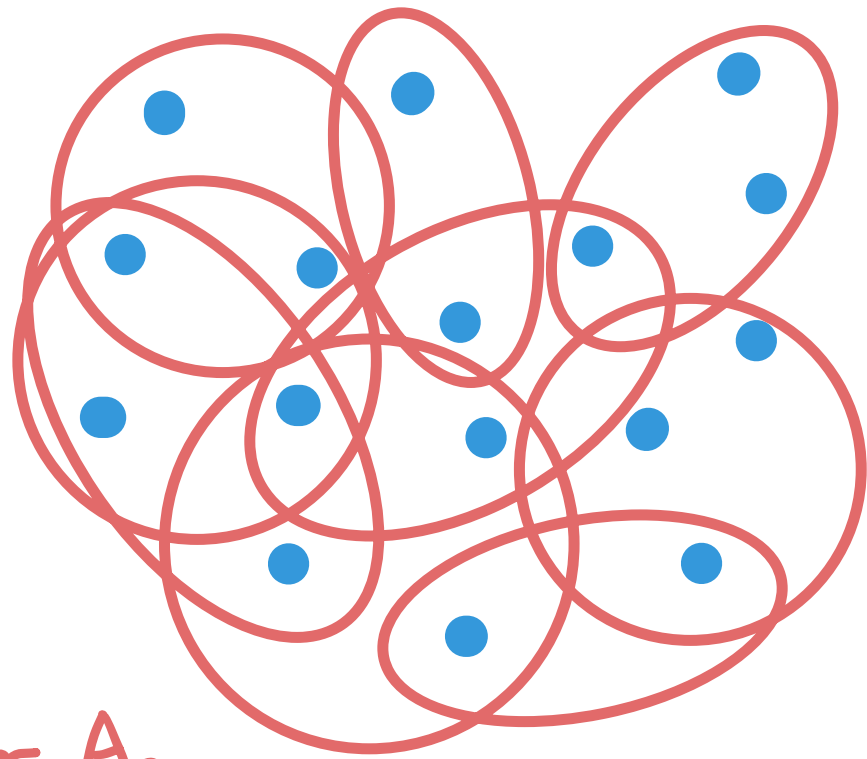
Randomized Greedy

Maximum Coverage

m points $[m] \equiv \{1, \dots, m\}$

n sets $A_1, \dots, A_n \subseteq [m]$

cardinality constraint $k \in \mathbb{N}$



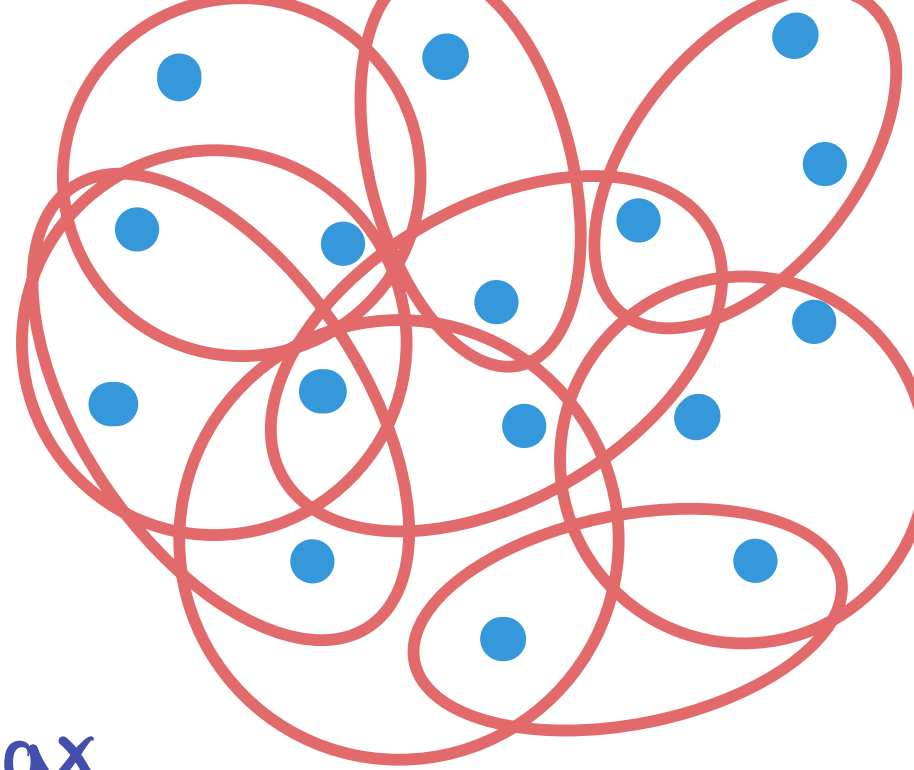
Goal: select k sets $A_{e_1}, A_{e_2}, \dots, A_{e_k}$

maximizing the union $|A_{e_1} \cup A_{e_2} \cup \dots \cup A_{e_k}|$

NP-Hard via set cover

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Set Cover
(next time)

minimize # sets
to cover all points

NP-Hard (via SAT)

LP + randomized
rounding

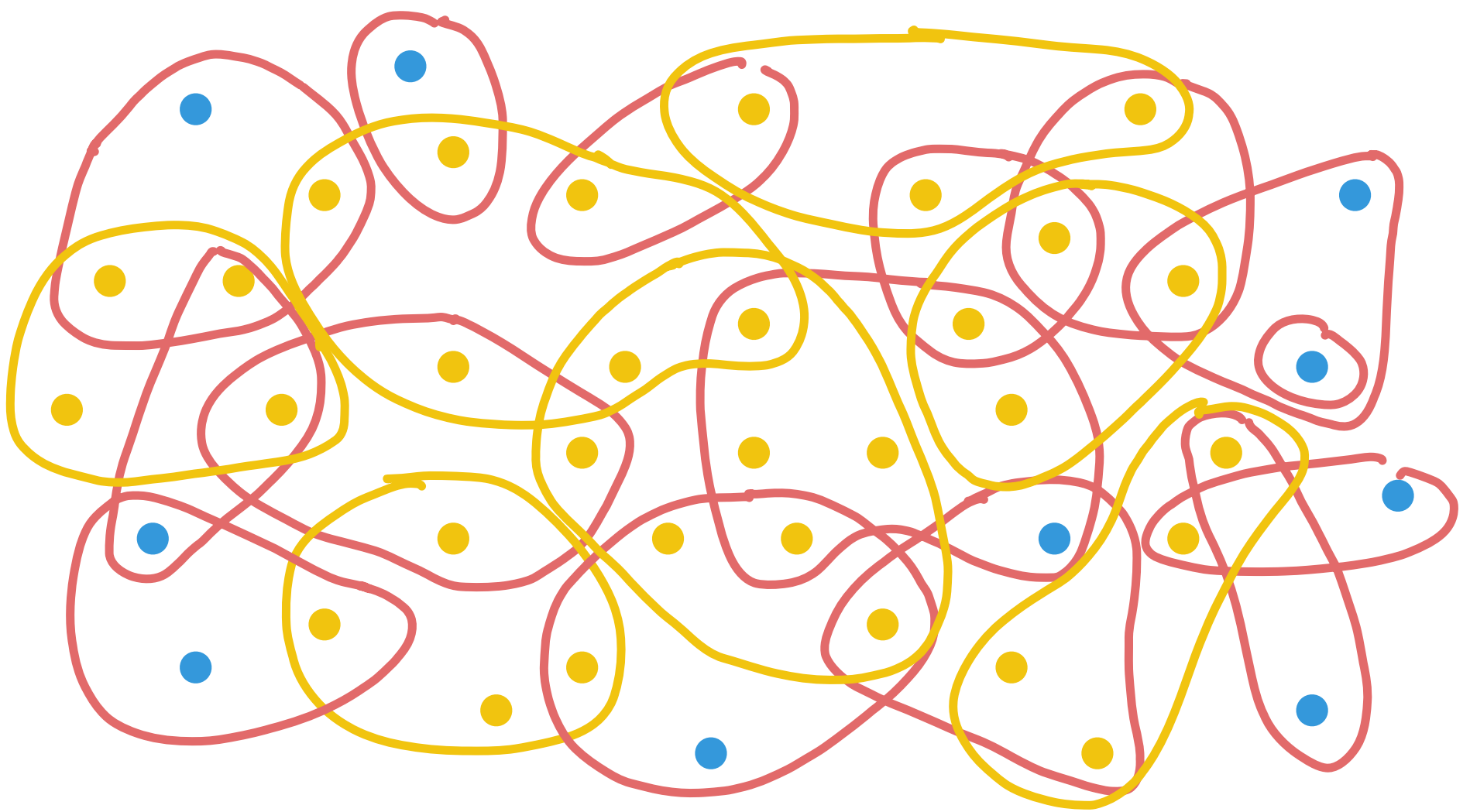
Max
Coverage

maximize # points
covered w/ $\leq k$ sets
(packing constraint)

(why?)
 $\Rightarrow$

NP-Hard

greedy algorithm



greedy algorithm repeatedly selects the

set covering the most new points

how good is the greedy algorithm?

(approximation ratio)

Set function $f: 2^N \rightarrow \mathbb{R}_{\geq 0}$ s.t.

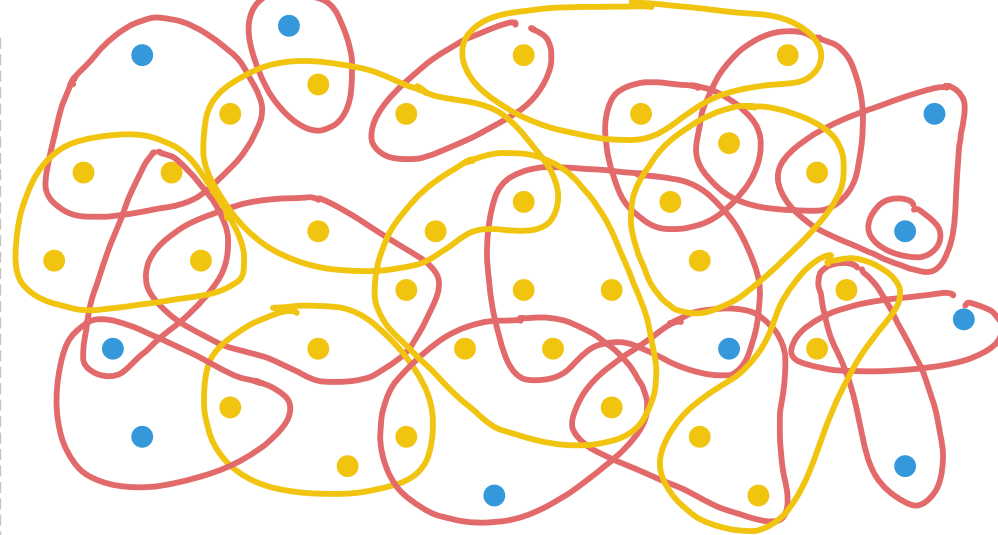
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- Monotone: $f(S) \leq f(T)$ for $S \subseteq T$
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equivalently:

"decreasing marginal values"

for $S \subseteq T \subseteq N$, $e \in N$,

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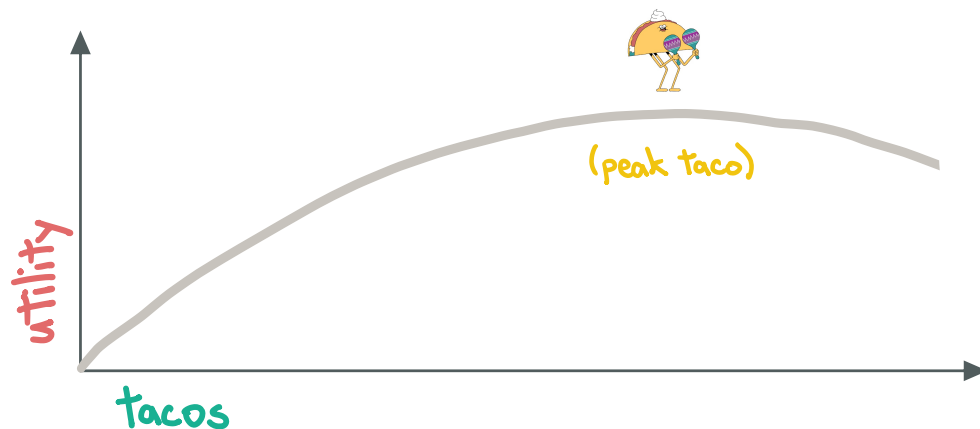


m points $[m] \equiv \{1, \dots, m\}$

n sets $A_1, \dots, A_n \subseteq [m]$

$S \subseteq [n]$ interpreted as
collection of sets (via indices)

$f(S) \stackrel{\text{def}}{=} \begin{pmatrix} \# \text{ points covered by} \\ \text{sets (w/ index in) } S \end{pmatrix}$



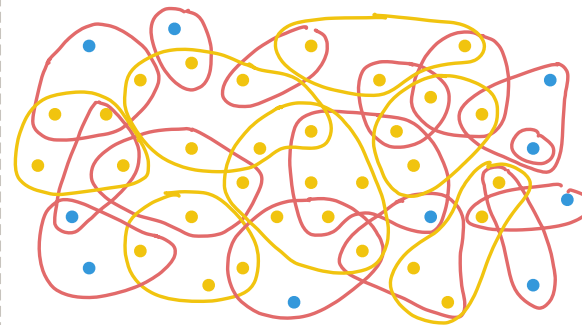
greedy($f, \mathcal{N}, k$)

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 - B. $S_i \leftarrow S_{i-1} + e_i$.
3. Return S_k .

Lemma $f(e_i \mid S_{i-1}) \geq \frac{1}{k} (\text{OPT} - f(S_{i-1}))$

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Suppose f was not monotone
does greedy work?

$$\begin{aligned} \leadsto f(S^* | S_i) &= f(S^* \cup S_i) - f(S_i) \\ &\stackrel{\text{monotonicity}}{\geq} f(S^*) - f(S_i) \end{aligned}$$

f not monotone: $f(S^* \cup S_i) \neq f(S^*)$

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Lemma $R \subseteq N$ random, $P[e \in R] \leq p \ \forall e$
then $E[f(R)] \geq (1-p)f(\emptyset)$

e.g. if S_i random, $P[e \in S_i] \leq 1/2 \ \forall e$

$$E[f(S^* \cup S_i)] \geq \frac{1}{2} f(S^*)$$

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Continuous extensions of f

$F: [0,1]^N \rightarrow \mathbb{R}$ extends $f: 2^N \rightarrow \mathbb{R}_{\geq 0}$
if $F(\mathbf{1}_S) = f(S) \quad \forall S \subseteq N$

Convex closure F^-

$F^-(p) \stackrel{\text{def}}{=} \inf \{E[f(R)] \text{ s.t. } P[e \in R] = p \quad \forall e\}$

Lovász:

$F^-(p) = E[f(S_t)] \quad \text{w/ } S_t = \{e: p_e \geq t\},$
 $t \in [0,1] \text{ unif. random}$

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Lovász: $F^-(p) = E[f(S_t)]$
w/ $S_t = \{e : p_e \geq t\}$, $t \in [0, 1]$ unif. random

proof:

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